Constructing random invertible matrices over a finite field

Sebastià Xambó-Descamps and Narcís Sayols Baixeras

Abstract

The aim of this note is to describe an iterative construction of a random invertible matrix of any positive order over a finite field, to show that the disbribution of this matrix is uniform, and to describe a PYECC implementation of the algorithm. The algorithm is inspired in the D. Randall report [1].

Notations and conventions

 $F=F_q$ will denote a finite field of cardinal q. The ring of F-matrices of order k will be denoted F(k) and the multiplicative group of the $A\in F(k)$ that are invertible by GL(k).

Given a random nonzero vector $v \in F^k$ we define the matrix I_v as the result of replacing the rth row of the identity matrix $I_k \in Gl(k)$ by v, where r is the index of the first non-zero component of v. It is clear that $I_v \in GL(k+1)$, as its determinant is equal to $v[r] \neq 0$.

1 Random extension of an invertible matrix

The main tool of this note consists of the following procedure:

- **1.1** (Procedure rd extend(A)). Let $A \in GL(k)$.
 - 1. Chose a random nonzero vector $v \in F^{k+1}$, and let r be the index of its first nonzero component.
 - 2. Let $e_r \in F^{k+1}$ be the vector whose entries are all 0, except for a 1 in the component of index r, and define $A' \in GL(k+1)$ so that $A'[0] = e_r$, $A'_{0,r} = A$, and with random values A'[j,r] drawn from F for $j=1,\ldots,k$.
 - 3. Return the matrix $B = A'I_v \in GL(k+1)$.
- **1.2** (Main argument). Let G_v be the subset of the $B \in GL(k+1)$ whose first row is a given non-zero vector v. Then it is immediate that the map $G_{e_r} \to G_v$, $A' \mapsto A'I_v$, is bijective (the inverse map is given by $B \mapsto BI_v^{-1} = BI_{\bar{v}}$, where $\bar{v}_j = 0$ for j < r, $\bar{v}_r = 1/v_r$, and $\bar{v}_j = -v_j/v_r$ for j < r). This shows that p(B) = p(v)p(A') = p(v)p(a)p(A), where $a = [a_{1r}, \ldots, a_{kr}]$, and the uniformity claim is clear because this value is a constant. Note that $p(v) = 1/(q^{k+1}-1)$, $p(a) = 1/q^k$, and $p(A) = 1/N_k$, where N_k is the cardinal of GL(k). If fact we have $p(B) = 1/N_{k+1}$, as it could not be otherwise, for it is easy to count that $N_k = (q^k 1) \cdots (q^k q^{k-1})$ and then $(q^{k+1} 1)q^kN_k = N_{k+1}$.
- **1.3** (The function rd_insert(A,r)). This function implements the construction of A', once r is known:

```
def rd_insert(A,r):
k = ncols(A)
if r<0 or r>n: return "r has to be in 0..k"
F = K_{-}(A)
A1 = matrix(F, k+1, k+1)
A1[0,r] = 1
for j in range(1,n+1):
    A1[j,r] = rd(F)
if r==0:
    A1[1:,1:] = A[:,:]
elif r==n:
    A1[1:,0:n] = A[:,:]
else:
   A1[1:,0:r] = A[:,0:r]
   A1[1:,r+1:n+1] = A[:,r:n]
return A1
```

1.4 (The function rd_extend(A)). Now we can write a function that implements the procedure rd extend:

```
def rd_extend(A):
k = ncols(A); F = K_(A)
v = rd_nonzero_vector(F,k+1)
r = 0
for j in range(k+1):
    if v[j]!=0:
        r = j; break
B = rd_insert(A,r)
x = v[r]
for j in range(r+1,k+1):
    B[:,j] = B[:,j]+ A[:,r]*v[j]
B[:,r] = x*A[:,r]
return B
```

2 The iterative procedure

Now the iterative procedure rd GL(n,F) can be obtained with the following function:

```
def rd_GL(n,F=Zn(2)):
a = rd_nonzero(F)
A = matrix([[a]])
for _ in range(2,n+1):
    A = rd_extend(A)
return A
```

Note that the first step guarantees that a is a chosen uniformly at random among the non-zero elements of F. Then we iterate rd_extend (n-1 times), and the main argument shows that at the end we get, by induction, a matrix of GL(n,F) chosen uniformly at random.

References

[1] D. Randall, "Efficient generation of random nonsingular matrices," 1991. Technical Report No. UCB/CSD-91-658, EECS Department, University of California, Berkeley.